# On social networks that support learning arXiv:2011.05255 

Itai Arieli, Fedor Sandomirskiy*, Rann Smorodinsky<br>*Technion, Haifa \& Higher School of Economics, St.Petersburg $\rightarrow$ Caltech<br>e-mail: fedor.sandomirskiy@gmail.com<br>homepage: https://www.fedors.info/

## Social learning

- each agent is going to make a single decision
- Android/iPhone, Private/Public kindergartens, restaurant A/B
- gets individual noisy signal about the best alternative \& observes choices made by predecessors


## Social learning

- each agent is going to make a single decision
- Android/iPhone, Private/Public kindergartens, restaurant A/B
- gets individual noisy signal about the best alternative \& observes choices made by predecessors
- usually: failure of information aggregation (herding)
- first agents take the wrong action $\Rightarrow$ others repeat it \& ignore their private signals $\Rightarrow$ information cascade (Banerjee [1992], Bikhchandani et al. [1992])


## Social learning

- each agent is going to make a single decision
- Android/iPhone, Private/Public kindergartens, restaurant A/B
- gets individual noisy signal about the best alternative \& observes choices made by predecessors
- usually: failure of information aggregation (herding)
- first agents take the wrong action $\Rightarrow$ others repeat it \& ignore their private signals $\Rightarrow$ information cascade (Banerjee [1992], Bikhchandani et al. [1992])
- mitigation measures
- signals of unbounded quality (Smith and Sorensen [2000])
- restricted observation: actions of friends on a social network (Smith [1991], Sgroi [2002], Acemoglu et al. [2010])


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade

The big puzzle
Which properties of the network are responsible for information aggregation?

- Are we the first to study this question?


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

- Are we the first to study this question? NO and YES
- NO
- topological conditions for a given ordering of agents Smith [1991], Sgroi [2002], Acemoglu et al. [2010]
- the timing of decisions determines social connections
- reasonable for life-long decisions (doctor/teacher) but not for (Android/iPhone)


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

- Are we the first to study this question? NO and YES
- NO
- topological conditions for a given ordering of agents Smith [1991], Sgroi [2002], Acemoglu et al. [2010]
- the timing of decisions determines social connections
- reasonable for life-long decisions (doctor/teacher) but not for (Android/iPhone)
- YES if the social structure and the timing of decisions are unrelated


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

- Are we the first to study this question? NO and YES
- NO
- topological conditions for a given ordering of agents Smith [1991], Sgroi [2002], Acemoglu et al. [2010]
- the timing of decisions determines social connections
- reasonable for life-long decisions (doctor/teacher) but not for (Android/iPhone)
- YES if the social structure and the timing of decisions are unrelated
- Our model: the network is given and the order is random


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

- Are we the first to study this question? NO and YES
- NO
- topological conditions for a given ordering of agents Smith [1991], Sgroi [2002], Acemoglu et al. [2010]
- the timing of decisions determines social connections
- reasonable for life-long decisions (doctor/teacher) but not for (Android/iPhone)
- YES if the social structure and the timing of decisions are unrelated
- Our model: the network is given and the order is random
- the network must aggregate information for most orders (very demanding!)


## Our question

- Agents are Bayesian-rational, sit on a network, act only once
- Signals have bounded quality
- cannot stop the information cascade


## The big puzzle

Which properties of the network are responsible for information aggregation?

- Are we the first to study this question? NO and YES
- NO
- topological conditions for a given ordering of agents Smith [1991], Sgroi [2002], Acemoglu et al. [2010]
- the timing of decisions determines social connections
- reasonable for life-long decisions (doctor/teacher) but not for (Android/iPhone)
- YES if the social structure and the timing of decisions are unrelated
- Our model: the network is given and the order is random
- the network must aggregate information for most orders (very demanding!)
- an example of such a network (Bahar et al. [2020])


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away
- no global information cascades
- quality of agent's decision is determined by his small neighborhood


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away $\Rightarrow$
- no global information cascades
- quality of agent's decision is determined by his small neighborhood


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away $\Rightarrow$
- no global information cascades
- quality of agent's decision is determined by his small neighborhood


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away $\Rightarrow$
- no global information cascades
- quality of agent's decision is determined by his small neighborhood
- Local learning requirement: the condition on agent's neighborhood for high-quality decision
- Want well-informed decisions? Make sure to be a part of mutually exclusive social circles!


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away $\Rightarrow$
- no global information cascades
- quality of agent's decision is determined by his small neighborhood
- Local learning requirement: the condition on agent's neighborhood for high-quality decision
- Want well-informed decisions? Make sure to be a part of mutually exclusive social circles!


## What will we see?

- Localization phenomenon: agent's decision is almost independent from those who are far away $\Rightarrow$
- no global information cascades
- quality of agent's decision is determined by his small neighborhood
- Local learning requirement: the condition on agent's neighborhood for high-quality decision
- Want well-informed decisions? Make sure to be a part of mutually exclusive social circles!
- Applications: constructing networks where learning is robust to disruptions


## The model and examples

## The model

The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.
- A finite Bayesian game $\Rightarrow$ an equilibrium exists


## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.
- A finite Bayesian game $\Rightarrow$ an equilibrium exists
- Learning qualities of an agent / of the network:

$$
\mathbb{P}\left(a_{v}=\theta\right) \quad / \quad L(G)=\frac{1}{|V|} \sum_{v \in V} \mathbb{P}\left(a_{v}=\theta\right)
$$

## The model

## The model

- undirected finite network $G=(V, E)$, vertices $=$ agents
- unobservable state $\theta \in\{$ blue, red $\}$ equally likely
- agent $v \in V$ arrives at $t_{v}$, i.i.d. uniform on $[0,1]$
- $v$ takes an action $a_{v} \in\{$ blue, red $\}$ depending on his information
- a binary signal that matches $\theta$ w.p. $p>\frac{1}{2}$ (i.i.d. conditional on $\theta$ )
- the set of friends who arrived earlier
- their actions
- the utility is 1 if $a_{v}=\theta$ and 0 , otherwise.
- A finite Bayesian game $\Rightarrow$ an equilibrium exists
- Learning qualities of an agent / of the network:

$$
\mathbb{P}\left(a_{v}=\theta\right) \quad / \quad L(G)=\frac{1}{|V|} \sum_{v \in V} \mathbb{P}\left(a_{v}=\theta\right)
$$

- the network supports learning if $L(G) \approx 1$


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals - 3rd agent repeats their action and ignores his signal - and so an


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals - 3rd agent repeats their action and ignores his signal - and so on


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals - 3rd agent repeats their action and ignores his signal - and so an


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals - 3rd agent repeats their action and ignores his signal - and so on


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals - 3rd agent repeats their action and ignores his signal - and so an


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals
- 3rd agent repeats their action and ignores his signal
- and so on


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals
- 3rd agent repeats their action and ignores his signal
- and so on


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals
- 3rd agent repeats their action and ignores his signal
- and so on


## Example: information cascade on a clique



- 1st and 2nd agents get the same blue signals
- 3rd agent repeats their action and ignores his signal
- and so on

The 1st two agents got wrong signals w.p. $(1-p)^{2} \Rightarrow$

$$
L\left(K_{n}\right) \leq 1-(1-p)^{2}
$$

## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.

## Theorem (Bahar et al. [2020])

$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.
Theorem (Bahar et al. [2020])
$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.
$m$ celebrities


## Example: celebrity graphs (Bahar et al. [2020])

$n$ commoners and $m$ celebrities observing each other: $n \gg m \gg 1$.

## Theorem (Bahar et al. [2020])

$\forall \delta>0$ there is a celebrity graph with $L \geq 1-\delta$.

$$
m \text { celebrities }
$$



## Remarks:

- the only known family of graphs with $L$ close to 1
- non-robustness: minority of celebrities is critical for learning


## Our results

## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $v$ observes $u$, i.e., $v u \in E$ and $t_{v}>t_{u}$


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $v$ observes $u$, i.e., $v u \in E$ and $t_{v}>t_{u}$
- $v$ observes $v_{1}$ who observes $u$


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $v$ observes $u$, i.e., $v u \in E$ and $t_{v}>t_{u}$
- $v$ observes $v_{1}$ who observes $u$
- $v$ observes $v_{1}$ who observes $v_{2}$ who observes $u$
- ...


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $v$ observes $u$, i.e., $v u \in E$ and $t_{v}>t_{u}$
- $v$ observes $v_{1}$ who observes $u$
- $v$ observes $v_{1}$ who observes $v_{2}$ who observes $u$
- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.

## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.
Proof: Show that no path ( $v \rightarrow$ boundary of $r$-neighborhood) is realized

- a path of length $r$ is realized with probability $1 /(r+1)$ !
- \#\{paths of length $r\} \leq D^{r}$
- the union bound


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.
Proof: Show that no path ( $v \rightarrow$ boundary of $r$-neighborhood) is realized

- a path of length $r$ is realized with probability $1 /(r+1)$ !
- the union bound


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.
Proof: Show that no path ( $v \rightarrow$ boundary of $r$-neighborhood) is realized

- a path of length $r$ is realized with probability $1 /(r+1)$ !
- \# \{paths of length $r\} \leq D^{r}$
- the union bound


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.
Proof: Show that no path ( $v \rightarrow$ boundary of $r$-neighborhood) is realized

- a path of length $r$ is realized with probability $1 /(r+1)$ !
- \#\{paths of length $r\} \leq D^{r}$
- the union bound


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.

## Implications:

- $d(v, u) \gg e \cdot D \Longrightarrow \mathbb{P}\left(a_{v}\right.$ and $a_{u}$ are dependent $)$ is exp. small
- impossibility of global information cascades
- the quality of $v$ 's decision is determined by the local structure of the


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.

## Implications:

- $d(v, u) \gg e \cdot D \Longrightarrow \mathbb{P}\left(a_{v}\right.$ and $a_{u}$ are dependent $)$ is exp. small


## - impossibility of global information cascades

- the quality of $v$ 's decision is determined by the local structure of the


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.

## Implications:

- $d(v, u) \gg e \cdot D \Longrightarrow \mathbb{P}\left(a_{v}\right.$ and $a_{u}$ are dependent $)$ is exp. small
- impossibility of global information cascades
- the quality of $v$ 's decision is determined by the local structure of the


## Localization phenomenon

Question: When can the action of $u$ affect the action of $v$ ?

- $\exists$ a path $\left(v=v_{0}, v_{1}, \ldots, v_{n-1}, v_{n}=u\right)$ such that $t_{v_{i}}>t_{v_{i+1}} \forall i$
- call such a path realized

Definition: Realized component

$$
N^{\text {real }}(v)=\{u: \exists \text { a realized path }(v \rightarrow u)\}
$$

## Proposition

$$
\mathbb{P}\left(N^{\text {real }}(v) \subset r \text {-neighborhood of } v\right) \geq 1-2\left(\frac{e \cdot D}{r}\right)^{r}
$$

where $D$ is the maximal degree.

## Implications:

- $d(v, u) \gg e \cdot D \Longrightarrow \mathbb{P}\left(a_{v}\right.$ and $a_{u}$ are dependent $)$ is exp. small
- impossibility of global information cascades
- the quality of $v$ 's decision is determined by the local structure of the network around $v$


## Local Learning Requirement



LLR with parameters $(d, r, D)$ :

- $v$ has a subset of $\geq d$ friends s.t.
- each of them has degree $\geq d$
- their $r$-neighborhoods in $G \backslash v$ are disjoint
- the max degree in these neighborhoods $\leq D$

$$
(d, r, D)=(3,2,7)
$$

## Local Learning Requirement



LLR with parameters ( $d, r, D$ ):

- $v$ has a subset of $\geq d$ friends s.t.
- each of them has degree $\geq d$
- their $r$-neighborhoods in $G \backslash v$ are disjoint
- the max degree in these neighborhoods $\leq D$

$$
(d, r, D)=(3,2,7)
$$

Theorem
$\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad$ where $\quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}$

## Local Learning Requirement



LLR with parameters $(d, r, D)$ :

- $v$ has a subset of $\geq d$ friends s.t.
- each of them has degree $\geq d$
- their $r$-neighborhoods in $G \backslash v$ are disjoint
- the max degree in these neighborhoods $\leq D$

$$
(d, r, D)=(3,2,7)
$$

Theorem
$\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad$ where $\quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}$

Proof:

## Local Learning Requirement



$$
(d, r, D)=(3,2,7)
$$

Theorem

$$
\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad \text { where } \quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}
$$

Proof:

- localization $\Rightarrow$ realized components of $v$ 's friends are disjoint with probability $\geq 1-\psi \Rightarrow$ independence


## Local Learning Requirement



$$
(d, r, D)=(3,2,7)
$$

Theorem

$$
\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad \text { where } \quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}
$$

Proof:

- localization $\Rightarrow$ realized components of $v$ 's friends are disjoint with probability $\geq 1-\psi \Rightarrow$ independence
- each friend takes correct action with prob. $\geq p \Rightarrow$ informativeness


## Local Learning Requirement



$$
(d, r, D)=(3,2,7)
$$

Theorem

$$
\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad \text { where } \quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}
$$

Proof:

- localization $\Rightarrow$ realized components of $v$ 's friends are disjoint with probability $\geq 1-\psi \Rightarrow$ independence
- each friend takes correct action with prob. $\geq p \Rightarrow$ informativeness
- $v$ observes $O(d)$ independent sources $\Rightarrow$ use Chernoff's bound. $\square$


## Local Learning Requirement



LLR with parameters ( $d, r, D$ ):

- $v$ has a subset of $\geq d$ friends s.t.
- each of them has degree $\geq d$
- their $r$-neighborhoods in $G \backslash v$ are disjoint
- the max degree in these neighborhoods $\leq D$

$$
(d, r, D)=(3,2,7)
$$

Theorem
$\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad$ where $\quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}$

## Local Learning Requirement



## Theorem

$$
\mathbb{P}\left(a_{v}=\theta\right) \geq 1-\left(\psi+\frac{18}{\sqrt{d-1}(2 p-1-\psi)}\right), \quad \text { where } \quad \psi=r \cdot\left(\frac{e \cdot D}{r}\right)^{r}
$$

Global implications of LLR: apply to each agent in the network

## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta
$$

## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Why surprising?
> - theory of the two-step information flow (Katz and Lazarsfeld [1955]): $\exists$ a minority critical for information-aggregation and predetermined by the network structure (opinion leaders)

- Bayesian social learning is fragile Frick et al. [2020], Mueller-Frank [2018], Bohren [2016]


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Why surprising?

- theory of the two-step information flow (Katz and Lazarsfeld [1955]): $\exists$ a minority critical for information-aggregation and predetermined by the network structure (opinion leaders)
- Bayesian social learning is fragile Frick et al. [2020], Mueller-Frank [2018], Bohren [2016]


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Why surprising?

- theory of the two-step information flow (Katz and Lazarsfeld [1955]): $\exists$ a minority critical for information-aggregation and predetermined by the network structure (opinion leaders)
- e.g., celebrities in Bahar et al. [2020]: if eliminated $\Rightarrow$ no aggregation



## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|} .
$$

Why surprising?

- theory of the two-step information flow (Katz and Lazarsfeld [1955]): $\exists$ a minority critical for information-aggregation and predetermined by the network structure (opinion leaders)
- e.g., celebrities in Bahar et al. [2020]: if eliminated $\Rightarrow$ no aggregation



## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|} .
$$

Why surprising?

- theory of the two-step information flow (Katz and Lazarsfeld [1955]): $\exists$ a minority critical for information-aggregation and predetermined by the network structure (opinion leaders)
- e.g., celebrities in Bahar et al. [2020]: if eliminated $\Rightarrow$ no aggregation
- Bayesian social learning is fragile Frick et al. [2020], Mueller-Frank [2018], Bohren [2016]


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|} .
$$

Proof:

- Apply LLR to each agent in $G$
- existence of such $G \Longleftarrow$ theory of expanders more datails


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|} .
$$

Proof:

- Apply LLR to each agent in $G$
- ( $\star$ ): $G$ is symmetric, high degrees, no short cycles
- ( $\star \star$ ): additionally, most $u \in U$ have high degrees in $\left.G\right|_{u}$ for $U \subset V$
- existence of such $G \Longleftarrow$ theory of expanders merd datis $\square$


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|}
$$

## Proof:

- Apply LLR to each agent in $G$
- ( $\star$ ): $G$ is symmetric, high degrees, no short cycles
- existence of such $G \Longleftarrow$ theory of expanders


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|} .
$$

## Proof:

- Apply LLR to each agent in $G$
- ( $\star$ ): $G$ is symmetric, high degrees, no short cycles
- ( $\star \star$ ): additionally, most $u \in U$ have high degrees in $\left.G\right|_{u}$ for $U \subset V$
- existence of such $G \Longleftarrow$ theory of expanders


## Applications: egalitarian societies and robust learning

Symmetry: $G=(V, E)$ is symmetric if for any $v, v^{\prime} \in V$, there is an automorphism $f$ such that $f(v)=v^{\prime}$.

## Proposition

For any $\delta>0$ there exists a symmetric network $G=(V, E)$ with

$$
L(G) \geq 1-\delta .
$$

Moreover, for any $U \subset V$, the sub-network satisfies

$$
L\left(\left.G\right|_{U}\right) \geq 1-\frac{\delta}{\alpha^{3}}, \quad \text { where } \quad \alpha=\frac{|U|}{|V|}
$$

## Proof:

- Apply LLR to each agent in $G$
- ( $\star$ ): $G$ is symmetric, high degrees, no short cycles
- ( $\star \star$ ): additionally, most $u \in U$ have high degrees in $\left.G\right|_{U}$ for $U \subset V$
- existence of such $G \Longleftarrow$ theory of expanders more deatils


## Summary

- Decoupling the network and the order of actions
- long paths of information transmission \& global cascades are unlikely
- learning quality of an agent is determined by the local structure
- LLR: a necessary condition for high quality \& no local cascades
- Bayesian models do not have explicit solutions
- Our approach is indirect. No insights in how equilibria look like.
- Future:
- How do equilibria look like? a simple open problem
- Other necessary and sufficient conditions for high learning quality


## Summary

- Decoupling the network and the order of actions
- long paths of information transmission \& global cascades are unlikely
- learning quality of an agent is determined by the local structure
- LLR: a necessary condition for high quality \& no local cascades
- Bayesian models do not have explicit solutions
- Our approach is indirect. No insights in how equilibria look like.
- Future:
- How do equilibria look like? a simple open problem
- Other necessary and sufficient conditions for high learning quality
Thank you!


## References

Daron Acemoglu, Munther A. Dahleh, Ilan Lobel, and Asuman Ozdaglar. Bayesian learning in social networks. Review of Economic Studies, 78: 1-34, 2010.
Noga Alon and Fan RK Chung. Explicit construction of linear sized tolerant networks. Discrete Mathematics, 72(1-3):15-19, 1988.
Gal Bahar, Itai Arieli, Rann Smorodinsky, and Moshe Tennenholtz. Multi-issue social learning. Mathematical Social Sciences, 104:29-39, 2020.

Abhijit V Banerjee. A simple model of herd behavior. The quarterly journal of economics, 107(3):797-817, 1992.
S. Bikhchandani, D. Hirshleifer, and I. Welch. A theory of fads, fashion, custom and cultural change as information cascade. The Journal of Political Economy, 100:992-1026, 1992.
J Aislinn Bohren. Informational herding with model misspecification. Journal of Economic Theory, 163:222-247, 2016.

Xavier Dahan. Regular graphs of large girth and arbitrary degree.
Combinatorica, 34(4):407-426, 2014.

Mira Frick, Ryota lijima, and Yuhta Ishii. Misinterpreting others and the fragility of social learning. Econometrica (forthcoming), 2020.

Elihu Katz and Paul F Lazarsfeld. Personal influence: the part played by people in the flow of mass communications. 1955.

Alexander Lubotzky, Ralph Phillips, and Peter Sarnak. Ramanujan graphs. Combinatorica, 8(3):261-277, 1988.

Manuel Mueller-Frank. Manipulating opinions in social networks. Available at SSRN 3080219, 2018.

Daniel Sgroi. Optimizing information in the herd: Guinea pigs, profits, and welfare. Games and Economic Behavior, 39:137-166, 2002.
L. Smith and P. Sorensen. Pathalogical outcomes of observational learning. Econometrica, 68:371-398, 2000.

Lones A Smith. Essays on dynamic models of equilibrium and learning. PhD thesis, University of Chicago, Department of Economics, 1991.

## Why expander graphs?

We need $G=(V, E)$ such that:

- ( $\star$ ) symmetric \& minimal degree is high \& no short cycles
- ( $\star \star$ ) most $u \in U$ have high degrees in $\left.G\right|_{U}, \forall U \subset V$ big enough

Definition: $d$-regular graph $G$ is an expander if $\lambda_{2}(G) \ll \lambda_{1}(G)=d$.

- the simple random walk forgets the origin fast
- $\Rightarrow$ best expanders have no short cycles and are highly connected


## Why expander graphs?

We need $G=(V, E)$ such that:

- ( $\star$ ) symmetric \& minimal degree is high \& no short cycles
- ( $\star \star$ ) most $u \in U$ have high degrees in $\left.G\right|_{U}, \forall U \subset V$ big enough

Definition: $d$-regular graph $G$ is an expander if $\lambda_{2}(G) \ll \lambda_{1}(G)=d$.

- the simple random walk forgets the origin fast
- $\Rightarrow$ best expanders have no short cycles and are highly connected

Ramanujan expanders (Lubotzky et al. [1988], Dahan [2014])
$\forall d \geq 11$ and $\forall g \geq 0 \exists$ symmetric $G$ with cycles $\geq g$ and $\lambda_{2} \leq 2 \sqrt{d-1}$

## Why expander graphs?

We need $G=(V, E)$ such that:

- ( $\star$ ) symmetric \& minimal degree is high \& no short cycles
- ( $\star \star$ ) most $u \in U$ have high degrees in $\left.G\right|_{U}, \forall U \subset V$ big enough

Definition: $d$-regular graph $G$ is an expander if $\lambda_{2}(G) \ll \lambda_{1}(G)=d$.

- the simple random walk forgets the origin fast
- $\Rightarrow$ best expanders have no short cycles and are highly connected


## Ramanujan expanders (Lubotzky et al. [1988], Dahan [2014])

$\forall d \geq 11$ and $\forall g \geq 0 \exists$ symmetric $G$ with cycles $\geq g$ and $\lambda_{2} \leq 2 \sqrt{d-1}$
For $U, U^{\prime} \subset V$, denote $E\left(U, U^{\prime}\right)=\left\{e \in E: e\right.$ connects $U$ and $\left.U^{\prime}\right\}$.

## Why expander graphs?

We need $G=(V, E)$ such that:

- ( $\star$ ) symmetric \& minimal degree is high \& no short cycles
- ( $\star \star$ ) most $u \in U$ have high degrees in $G \mid U, \forall U \subset V$ big enough

Definition: $d$-regular graph $G$ is an expander if $\lambda_{2}(G) \ll \lambda_{1}(G)=d$.

- the simple random walk forgets the origin fast
- $\Rightarrow$ best expanders have no short cycles and are highly connected


## Ramanujan expanders (Lubotzky et al. [1988], Dahan [2014])

$\forall d \geq 11$ and $\forall g \geq 0 \exists$ symmetric $G$ with cycles $\geq g$ and $\lambda_{2} \leq 2 \sqrt{d-1}$
For $U, U^{\prime} \subset V$, denote $E\left(U, U^{\prime}\right)=\left\{e \in E\right.$ : e connects $U$ and $\left.U^{\prime}\right\}$.
Mixing lemma (Alon and Chung [1988])
$\left|E\left(U, U^{\prime}\right)\right|=\frac{d}{|V|} \cdot|U| \cdot\left|U^{\prime}\right|+\tau, \quad$ where $\quad|\tau| \leq \lambda_{2} \sqrt{|U|\left|U^{\prime}\right|}$.

## Why expander graphs?

We need $G=(V, E)$ such that:

- $(\star)$ symmetric \& minimal degree is high \& no short cycles
- ( $\star \star$ ) most $u \in U$ have high degrees in $\left.G\right|_{U}, \forall U \subset V$ big enough

Definition: $d$-regular graph $G$ is an expander if $\lambda_{2}(G) \ll \lambda_{1}(G)=d$.

- the simple random walk forgets the origin fast
- $\Rightarrow$ best expanders have no short cycles and are highly connected


## Ramanujan expanders (Lubotzky et al. [1988], Dahan [2014])

$\forall d \geq 11$ and $\forall g \geq 0 \exists$ symmetric $G$ with cycles $\geq g$ and $\lambda_{2} \leq 2 \sqrt{d-1}$
For $U, U^{\prime} \subset V$, denote $E\left(U, U^{\prime}\right)=\left\{e \in E\right.$ : e connects $U$ and $\left.U^{\prime}\right\}$.
Mixing lemma (Alon and Chung [1988])
$\left|E\left(U, U^{\prime}\right)\right|=\frac{d}{|V|} \cdot|U| \cdot\left|U^{\prime}\right|+\tau, \quad$ where $\quad|\tau| \leq \lambda_{2} \sqrt{|U|\left|U^{\prime}\right|}$.
$(\star \star)$ : if $|U|=\alpha|V|$, the average degree in $\left.G\right|_{U}$ is $\frac{|E(U, U)|}{|U|} \approx \alpha \cdot d$.

## Open problem: puzzling unanimity Gack osmmay

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\left.\mid a_{v}=r e d\right) \geq p$

## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\left.\mid a_{v}=r e d\right) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?

## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\operatorname{red} \mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?

Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?
Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

Difficulty:
non-monotonicity of the posterior: more red actions observed may signal about herding $\Rightarrow$ weaker evidence.

## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?
Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

## Difficulty:

non-monotonicity of the posterior: more red actions observed may signal about herding $\Rightarrow$ weaker evidence.


## Example with fixed arrival order

- Strong evidence for $\theta=$ red?
- What if one observation was blue?


## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?
Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

## Difficulty:

non-monotonicity of the posterior: more red actions observed may signal about herding $\Rightarrow$ weaker evidence.


## Example with fixed arrival order

- Strong evidence for $\theta=$ red? NO
- What if one observation was blue?


## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?
Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

## Difficulty:

non-monotonicity of the posterior: more red actions observed may signal about herding $\Rightarrow$ weaker evidence.


## Example with fixed arrival order

- Strong evidence for $\theta=$ red? NO
- What if one observation was blue?


## Open problem: puzzling unanimity

For any agent $v, \mathbb{P}\left(\theta=\right.$ red $\mid a_{v}=$ red $) \geq p$

## Question

Is this true for groups? Namely,
$\mathbb{P}\left(\theta=\right.$ red $\left.\mid\left(a_{v}\right)_{v \in U}=r e d\right) \geq p \quad$ for any $U \subset V$ ?
Remark: if yes, an agent observing $U$, will (weakly) prefer the unanimous decision to his own signal $\Rightarrow$ red propagates.

## Difficulty:

non-monotonicity of the posterior: more red actions observed may signal about herding $\Rightarrow$ weaker evidence.


## Example with fixed arrival order

- Strong evidence for $\theta=$ red? NO
- What if one observation was blue? then YES


## Robustness to random-subset elimination

## Proposition

Arbitrary $G=(V, E)$ with learning quality $L(G)=1-\delta$.
For uniformly random $U \subset V$ such that $|U|=\alpha \cdot|V|$, the subnetwork
satisfies

$$
\mathbb{E}[L(G \mid U)] \geq 1-\frac{\delta}{\alpha}
$$

## Robustness to random-subset elimination

## Proposition

Arbitrary $G=(V, E)$ with learning quality $L(G)=1-\delta$.
For uniformly random $U \subset V$ such that $|U|=\alpha \cdot|V|$, the subnetwork satisfies

$$
\mathbb{E}[L(G \mid U)] \geq 1-\frac{\delta}{\alpha}
$$

## Proof sketch

- Coupling between learning on $G$ and the choice of $U$ :

$$
U=\{\text { the set of } \alpha \cdot|V| \text { earliest arrivals }\} .
$$

- Learning on $G \mid u$ becomes a part of learning on $G \Rightarrow$

$$
L(G) \leq \alpha \cdot \mathbb{E}\left[L\left(\left.G\right|_{U}\right)\right]+(1-\alpha) \cdot 1
$$

## Robustness to random-subset elimination

## Proposition

Arbitrary $G=(V, E)$ with learning quality $L(G)=1-\delta$.
For uniformly random $U \subset V$ such that $|U|=\alpha \cdot|V|$, the subnetwork satisfies

$$
\mathbb{E}[L(G \mid U)] \geq 1-\frac{\delta}{\alpha}
$$

## Proof sketch

- Coupling between learning on $G$ and the choice of $U$ :

$$
U=\{\text { the set of } \alpha \cdot|V| \text { earliest arrivals }\} .
$$

- Learning on $\left.G\right|_{U}$ becomes a part of learning on $G \Rightarrow$

$$
L(G) \leq \alpha \cdot \mathbb{E}\left[L\left(\left.G\right|_{U}\right)\right]+(1-\alpha) \cdot 1 .
$$

## Robustness to random-subset elimination

## Proposition

Arbitrary $G=(V, E)$ with learning quality $L(G)=1-\delta$.
For uniformly random $U \subset V$ such that $|U|=\alpha \cdot|V|$, the subnetwork satisfies

$$
\mathbb{E}\left[L\left(\left.G\right|_{U}\right)\right] \geq 1-\frac{\delta}{\alpha}
$$

## Proof sketch

- Coupling between learning on $G$ and the choice of $U$ :

$$
U=\{\text { the set of } \alpha \cdot|V| \text { earliest arrivals }\} .
$$

- Learning on $\left.G\right|_{U}$ becomes a part of learning on $G \Rightarrow$

$$
L(G) \leq \alpha \cdot \mathbb{E}\left[L\left(\left.G\right|_{U}\right)\right]+(1-\alpha) \cdot 1 .
$$

Example: celebrity graph, $\alpha=50 \% \Rightarrow \simeq 50 \%$ celebrities remain.

